

# Bidirectional leveled enumerators

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# Enumeration : what for ?

Enumeration is essential

- ▶ In general

- ▶ for infinite sets of objects
- ▶ for sets that are too big to be represented *in extenso*  
Python : `for i in range(n):`

- ▶ In particular

- ▶ search problems with large answer sets
- ▶ queries on large databases : iterators in Java and SQL :
- ▶ constraint satisfaction problems

Enumeration strategies for solving constraint satisfaction problems : A performance evaluation.

*Advances in Intelligent Systems and Computing*, 347 :169–179, 07 2015.

- ▶ ...

# Framework

Development of the **TRAG** system : Term Rewriting Automata and Graphs

- ▶ <https://idurand@bitbucket.org/idurand/trag.git>
- ▶ core system entirely written in **Common Lisp**
- ▶ A web interface<sup>2</sup> : [trag.labri.fr](http://trag.labri.fr) (uses some JavaScript)
  
- ▶ need for enumeration  $\Rightarrow$  Enum package, self-contained
- ▶ presented at ELS2012 in Zadar

# Overview of the talk

- ▶ Definition of an **enumerator**
- ▶ Presentation of the Enum package
- ▶ Problems raised by the enumeration of products
- ▶ Bidirectional leveled enumerators for enumerating products
- ▶ Conclusion

## Definition of an enumerator

An **enumerator**  $E$  is a state machine which outputs the elements of a sequence  $\hat{E} = e_0, e_1, \dots$  one at a time. The sequence may be finite  $(e_n)_{n \in [0, c[}$  or infinite  $(e_n)_{n \in \mathbb{N}}$ .

Examples of sequences

1. the sequence of the days of the week :  
*monday, tuesday, wednesday, thursday, friday, saturday, sunday*
2. the sequence of all natural integers :  $0, 1, 2, \dots$
3. the alternating sequence :  $1, -1, 1, -1, \dots$
4. the sequence of prime numbers :  $2, 3, 5, 7, 11, 13, \dots$

Two basic operations :

- ▶ is there a next element? (**next-element-p**)
- ▶ output the next element and move on (**next-element**)

## API for Enumerator (basic operations)

```
(defclass abstract-enumerator () ())  
  
(defgeneric next-element-p (enumerator)  
  (:documentation "returns NIL if there is no next element,  
  a non NIL value otherwise"))  
  
(defgeneric next-element (enumerator)  
  (:documentation "returns the next element,  
  moves to the following one"))
```

# Enumerator (main operation)

```
(defgeneric call-enumerator (enumerator)
  (:documentation
   "return as first value the next element of ENUMERATOR
    if it exists NIL otherwise
    as second value T if element was produced")
  (:method ((e abstract-enumerator))
    (if (next-element-p e)
        (values (next-element e) t)
        (values nil nil))))
```

## Example of the list enumerator

```
ENUM> (setq *abc* (make-list-enumerator '(a b c)))  
=> #<LIST-ENUMERATOR {100ADDAAC3}>  
ENUM> (next-element *abc*) => A  
ENUM> (next-element-p *abc*) => T  
ENUM> (next-element *abc*) => B  
ENUM> (call-enumerator *abc*)  
C  
T  
ENUM> (call-enumerator *abc*)  
NIL  
NIL  
ENUM> (collect-enum *abc*) => (A B C) ; only if finite  
ENUM> (defparameter *ab-i* (make-list-enumerator '(a b) :circ t))  
*I*  
ENUM> (collect-n-enum *i* 10)  
(A B A B A B A B A B)
```



# Simple vs relying enumerators

- ▶ **simple** enumerators : do not rely on other enumerators
  - ▶ constant enumerator
  - ▶ enumerator of the elements of Lisp sequence (`list`, `vector`)
  - ▶ enumerator of a sequence defined inductively
  - ▶ ...
- ▶ **relying** enumerators : rely on other enumerators
  - ▶ `sequential( $E^1, \dots, E^n$ )`
  - ▶ `parallel( $f, (E^1, \dots, E^n)$ )`
  - ▶ `filter( $p, E$ )`
  - ▶ `product( $f, (E^1, \dots, E^n)$ )`
  - ▶ ...

## A relying enumerator : parallel-enumerator

The parallel-enumerator is like the enumerator version of the `(mapcar function list &rest more-lists)`.

```
ENUM> (mapcar #'list
             '(e^1_0 e^1_1 e^1_2 e^1_3)
             '(e^2_0 e^2_1 e^2_2)
             '(e^3_0 e^3_1 e^3_2 e^3_3))
((E^1_0 E^2_0 E^3_0) (E^1_1 E^2_1 E^3_1) (E^1_2 E^2_2 E^3_2))
```

It stops with the **shortest** sequence.

Let  $E = \text{parallel}(f, (E^1, \dots, E^n))$ , each  $E^i$  enumerating  $e_0^i, e_1^i, \dots, e_{k_i}^i$ . Let  $k = \min(k_1, \dots, k_n)$ .

$$\begin{aligned} & f(e_0^1, e_0^2, \dots, e_0^{n-1}, e_0^n), \\ & f(e_1^1, e_1^2, \dots, e_1^{n-1}, e_1^n), \\ & \dots, \\ & f(e_k^1, e_k^2, \dots, e_k^{n-1}, e_k^n) \end{aligned}$$

# Injective enumerators

$$\begin{aligned}\widehat{E} : \mathbb{N} &\rightarrow D \\ n &\mapsto e_n\end{aligned}$$

A sequence  $\widehat{E}$  is an application from  $\mathbb{N}$  to some domain  $D$ .

Not necessarily **injective** : we may have  $e_i = e_j$  with  $i \neq j$ .

Example :  $e_0 = 1, e_1 = -1, e_2 = 1, e_3 = -1, \dots$

For simplifying the proofs, it is easier to assume that enumerators are injective.

This yields no loss of generality :

a non injective enumerator may always be transformed into an injective one by putting it in parallel with an enumerator of  $\mathbb{N}$ .

## Exemple of the parallel enumerator

```
ENUM> (setq *naturals* (make-inductive-enumerator 0 #'1+))
#<INDUCTIVE-ENUMERATOR {100C836033}>
ENUM> (collect-n-enum *naturals* 10) => (0 1 2 3 4 5 6 7 8 9)
ENUM> (setq *parallel*
        (make-parallel-enumerator (list *naturals* *abc*)))
#<PARALLEL-ENUMERATOR {10020EE5C3}>
ENUM> (collect-enum *parallel*) => ((0 A) (1 B) (2 C))
ENUM> (setq *parallel*
        (make-parallel-enumerator
         (list *naturals* (make-constant-enumerator 'a))))
#<PARALLEL-ENUMERATOR {10020F0893}>
ENUM> (collect-n-enum *parallel* 10)
((0 A) (1 A) (2 A) (3 A) (4 A) (5 A) (6 A) (7 A) (8 A) (9 A))
```

## Application : implementation of the `map` function

```
(defun map (result-type fun &rest sequences)
  (let ((enumerator
        (make-parallel-enumerator
         (mapcar
          (lambda (s)
            (make-sequence-enumerator s))
          sequences)
         :fun (lambda (tuple)
                (apply fun tuple))))))
    (if (null result-type)
        nil
        (coerce (collect-enum enumerator) result-type))))
```

```
ENUM> (map 'list #'list '(1 2 3) #(a b c d))
((1 A) (2 B) (3 C))
```

## Application : Erathostenes's sieve

```
(defclass erathostenes (unary-relying-enumerator)
  ((enum :type abstract-enumerator :accessor enum
         :initform (make-inductive-enumerator 2 #'1+))))

(defmethod next-element ((e erathostenes))
  (let ((prime (next-element (enum e))))
    (setf (enum e)
          (make-instance
           'filter-enumerator
           :enum (enum e) :fun (lambda (n) (plusp (mod n prime))))))
  prime))

(defun make-erathostenes-enumerator () (make-instance 'erathostenes))

ENUM> (setq *e* (make-erathostenes-enumerator))
#<ERATHOSTENES {100C268E03}>
ENUM> (collect-n-enum *e* 10)
(2 3 5 7 11 13 17 19 23 29)
ENUM> (collect-n-enum *e* 10 :init nil)
(31 37 41 43 47 53 59 61 67 71)
```

# Relying enumerators

$E^1, \dots, E^n$   $\rightarrow$  sequential  $\rightarrow$  sequential( $E^1, \dots, E^n$ )

$E^1, \dots, E^n$   
 $f$   $\rightarrow$  parallel  $\rightarrow$  parallel( $f, (E^1, \dots, E^n)$ )

$E$   $\rightarrow$  append  $\rightarrow$  append( $E$ )

$E$   
 $p$   $\rightarrow$  filter  $\rightarrow$  filter( $p, E$ )

$E$   
 $f$   
 $map$   $\rightarrow$  mapping  $\rightarrow$  mapping( $map, f, E$ )

$E^1, \dots, E^n$   
 $f$   $\rightarrow$  product  $\rightarrow$  product( $f, E^1, \dots, E^n$ )

# Product enumerators

Let  $E^1, \dots, E^p$  be nonempty **injective** enumerators s.t each  $E^i$  enumerates

- ▶  $e_0^i, e_1^i, \dots$  if  $E^i$  is infinite
- ▶  $e_0^i, e_1^i, \dots, e_{c^i-1}^i$  where  $c^i = \text{card}(E^i)$  otherwise.

Let  $T^p = \widehat{E}^1 \times \widehat{E}^2 \dots \times \widehat{E}^p$  be the **cartesian product** of the the  $\widehat{E}^i$ s.

$$T^p = \{(e_{j_1}^1, e_{j_2}^2, \dots, e_{j_p}^p) \mid \forall i \in [1, p], e_{j_i}^i \in E^i\}.$$

$E^i$ s injective  $\Rightarrow$  tuples  $(e_{j_1}^1, e_{j_2}^2, \dots, e_{j_p}^p)$  with distinct indices are distinct.

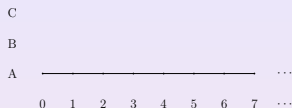
If every  $E^i$  is finite,  $\text{card}(T^p) = \prod_{i=1}^p c^i$  otherwise  $T^p$  is infinite.

**Multiple ways** of **ordering**  $T^p$  so multiple possible enumerators of  $T^p$ .

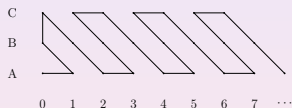


# Fairness property

Necessity of a **fair** ordering as soon as one of the  $E^i$  is infinite.



An unfair ordering  $\text{*naturals*} \times \text{*abc*}$



A fair diagonal ordering of  $\text{*naturals*} \times \text{*abc*}$

```
ENUM> (collect-n-enum
      (make-diagonal-product-enumerator *naturals* *abc*)
      20)
((0 A) (1 A) (0 B) (0 C) (1 B) (2 A) (3 A) (2 B) (1 C) (2 C)
 (3 B) (4 A) (5 A) (4 B) (3 C) (4 C) (5 B) (6 A) (7 A) (6 B))
```

This diagonal ordering was already there in 2012.

# Bidirectional enumerators

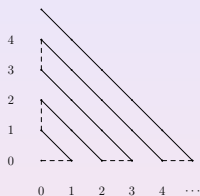
A **bidirectional** enumerator can move forward and backward.  
It has a **way** and operations to deal with it.

```
(defun next-element-p (B) (way-next-element-p (way B) B))
(defun next-element (B) (way-next-element (way B) B))

ENUM> (setq *b-naturals* (make-bidirectional-enumerator *naturals*))
#<BITIRECTIONAL-ENUMERATOR {100D46E113}>
ENUM> (way *b-naturals*) => 1
ENUM> (next-element *b-naturals*) => 0
ENUM> (next-element *b-naturals*) => 1
ENUM> (next-element *b-naturals*) => 2
ENUM> (next-element *b-naturals*) => 3
ENUM> (latest-element *b-naturals*) => 3
ENUM> (way-next-element -1 *b-naturals*) => 2
ENUM> (next-element *b-naturals*) => 3
ENUM> (invert-way *b-naturals*) => -1
ENUM> (next-element *b-naturals*) => 2
ENUM> (next-element *b-naturals*) => 1
ENUM> (way-next-element 1 *b-naturals*) => 2
ENUM> (next-element *b-naturals*) => 1
```

# Diagonal ordering of binary product : sliding and corner steps

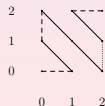
A pair of consecutive elements of an enumerator  $E$  is called a **step**.



A diagonal ordering of  $\mathbb{N} \times \mathbb{N}$

Dashed line : **sliding** step

With finite enumerators also **corner** steps (dotted line).



A diagonal ordering of  $[0, 2] \times [0, 2]$

## Distance between tuples

In an ordering of the cartesian product  $T^P$ , we define a **distance** between two enumerated tuples  $t_j = (e_{j^1}^1, \dots, e_{j^p}^p)$  and  $t_k = (e_{k^1}^1, \dots, e_{k^p}^p)$  as  $d(t_j, t_k) = \sum_{i=1}^p |k^i - j^i|$ .

To simplify, from now on we assume that **every  $E^i$  enumerates integers starting from 0**.

A list of **indices**  $(j^1, \dots, j^p)$  will be the same as the **tuple**  $(e_{j^1}^1, \dots, e_{j^p}^p) : \forall i \in [1; p], \forall j, e_j^i = j$

$$d((0 \ 0), (1 \ 0)) = 1, d((0 \ 2), (1 \ 1)) = 2$$

Each enumerator is like the **wheel** of a **mechanical counter** which can be moved forward or backward of one or more times.  
distance between two tuples = sum of **sizes of moves** of the wheels

## $d$ -orderings of a cartesian product $T^p$

**size** of a step  $(t_j, t_{j+1}) : d(t_j, t_{j+1})$ .

An ordering of  $T^p$  is a  **$d$ -ordering** if the size of each step is **atmost**  $d$ .

Our aim is to have a **2-ordering** of  $T^p$ ,  $\forall p \geq 2$ .

For  $p = 2$ , a diagonal ordering is a 2-ordering because each enumerator moves at most one notch forward or backward.

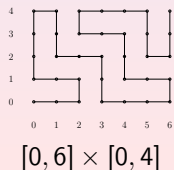
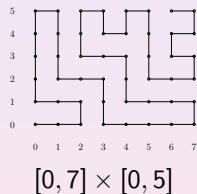
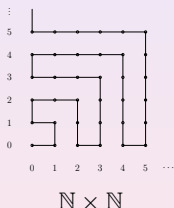
# 1-orderings for $p = 2$

There exist a 1-ordering for all binary products.

**not feasible** in our setting for arbitrary (finite or non finite) enumerators

need to know one step in advance whether we have reached the end of a finite enumerator to turn around before it is too late.

**next-element-p** not enough



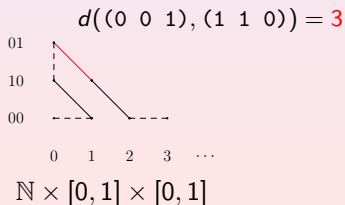
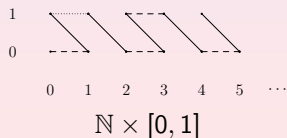
depends on the parity of the size of the finite enumerators

## 2-orderings

Our aim is to have a 2-ordering of  $T^P$ .

Recursive use of the binary diagonal  $\Rightarrow$   $p$ -ordering!

```
ENUM> (defparameter *e2* (make-list-enumerator '(0 1))) => *E2*
ENUM> (defparameter *p2*
      (make-diagonal-product-enumerator *naturals* *e2*)) => *P2*
ENUM> (collect-n-enum *p2* 11)
((0 0) (1 0) (0 1) (1 1) (2 0) (3 0) (2 1) (3 1) (4 0) (5 0) (4 1))
ENUM> (defparameter
      *p3*
      (make-diagonal-product-enumerator *e2* *p2* :fun #'cons))
*P3*
ENUM> (collect-n-enum *p3* 5)
((0 0 0) (1 0 0) (0 1 0) (0 0 1) (1 1 0))
```



## leveled orderings of $T^P$

**level** : subset of tuples with identical height.

**height** of a tuple  $t = (e_{j_1}^1, e_{j_2}^2, \dots, e_{j_p}^p) \in T^P$  is the sum of the indices of the elements in the  $\widehat{E}^i$  :  $h(t) = \sum_{i=1}^p j^i$

Note that with our hypothesis (each  $\widehat{E}^i$  is a prefix of  $\mathbb{N}$ ), the **height** of a tuple is the **sum** of its elements.

**$h^{\text{th}}$ -level** of  $T^P$  : set of tuples with height  $h$  (denoted by  $L_h$ ).

An ordering of  $T^P$  is **leveled** if it traverses the levels  $L_0, L_1, \dots$  in the **increasing** order of levels  $L_0, L_1, \dots$  (without constraint so far on the order of enumeration inside a level).



# Leveled enumerators

**leveled** enumerators follow a **leveled** ordering :  $L_0, L_1, \dots$

A step giving a change of level is called a **major step**.

A step inside a level is called a **minor step**.

leveled enumerators have predicate :

▶ `minor-step-p (E)`

which returns **T** if the next step (`next-element`) does not change the level (**NIL** otherwise).

In other words, it returns **false** when we are done with the enumeration of the current level.

# Bidirectional leveled enumerators

**bidirectional leveled** enumerator : leveled **and** bidirectional

**forward** : levels enumerated in **increasing** order :  $L_0, L_1, \dots$

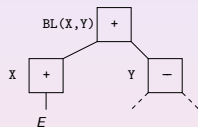
**backward** : levels enumerated in **decreasing** order :  $L_i, L_{i-1}, \dots$

```
ENUM> (defparameter *e3* (make-list-enumerator '(0 1 2))) => *e3*
ENUM> (defparameter
      *e* (make-bidirectional-leveled-enumerator *e3* *e3*)) => *E*
ENUM> (next-element *e*) => (0 0) ; level 0
ENUM> (next-element *e*) => (1 0) ; level 1
ENUM> (next-element *e*) => (0 1)
ENUM> (next-element *e*) => (0 2) ; level 2
ENUM> (next-element *e*) => (1 1)
ENUM> (next-element *e*) => (2 0)
ENUM> (next-element *e*) => (2 1) ; level 3
ENUM> (next-element *e*) => (1 2)
ENUM> (next-element *e*) => (2 2) ; level 4
ENUM> (invert-way *e*) => -1
ENUM> (next-element *e*) => (2 1) ; level 3
ENUM> (next-element *e*) => (1 2)
ENUM> (next-element *e*) => (0 2) ; level 2
```

# Product of bidirectional and bidirectional leveled

Let  $X$  be a **bidirectional** enumerator and  $Y$  be a **bidirectional leveled** enumerator,

We define  $D = BL(X, Y)$ , the **leveled bidirectional product** of  $X$  and  $Y$ .



When  $D = BL(X, Y)$  is created, the initial way of  $X$  is set to  $+1$  and the initial way of  $Y$  is set to  $-1$ .

The **minor** steps (same level) are **diagonal** steps

The **major** steps are either **sliding** or **corner** steps.

depending on the way of the enumerator the level **increases** or **decreases** by 1.

# Code for the product

```
(defun minor-step-p (D)
  ;; precondition (next-element-p D)
  (and (next-element-p (enum-y D))
        (or (next-element-p (enum-x D))
              (minor-step-p (enum-y D)))))

(defun way-next-element-p (way D)
  (or (way-next-element-p (way D) (enum-x D))
      (way-next-element-p (way D) (enum-y D))))

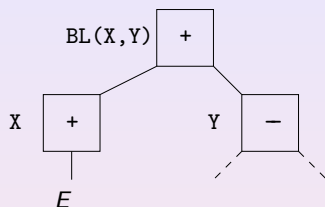
(defun sliding-step (X Y way)
  ;; precondition: X or Y can move in its way
  (if (next-element-p Y)
      (way-next-element way Y)
      (way-next-element way X))
  (invert-way X)
  (invert-way Y))

(defun corner-step (enum1 enum2 way)
  (when (plusp (* way (way enum1))))
    (psetf enum1 enum2 enum1))
  ;; enum1 is now the one that goes
  ;; in direction -way
  (invert-way enum1)
  ;; enum1 will move in direction way
  (next-element enum1)
  ;; enum2 will move in direction -way
  (invert-way enum2))

(defun way-next-element (way D)
  (let* ((enum-x (enum-x enum))
         (enum-y (enum-y enum))
         (next-x (next-element-p enum-x))
         (next-y (next-element-p enum-y)))
    (cond
     ((and next-y (minor-step-p enum-y))
      ;; lower-level minor step
      (next-element enum-y))
     ((and next-x next-y) ; minor-step on level
      ;; each one makes a major in its way
      (next-element enum-x) (next-element enum-y))
     ;; major step
     ((not (or next-x next-y))
      (corner-step enum-x enum-y way))
     (t (sliding-step enum-x enum-y way))))
  (latest-element enum))

(defun latest-element (D)
  (cons (latest-element(enum-x D))
        (latest-element (enum-y D))))
```

# Properties of the bidirectional leveled product



$BL(X, Y)$  is a bidirectional leveled enumerator.

**bidirectional** : by construction

**leveled** : by proof (see the paper)

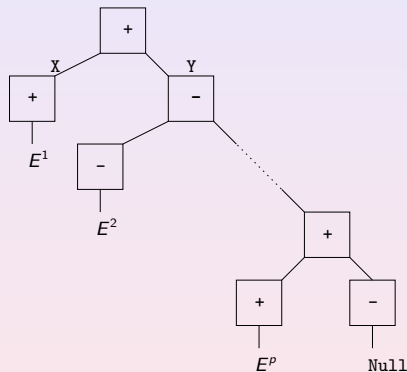
Moreover :

**Lemma** : if  $Y$  is a 2-ordering then so is  $BL(X, Y)$

**Proof** : see the paper

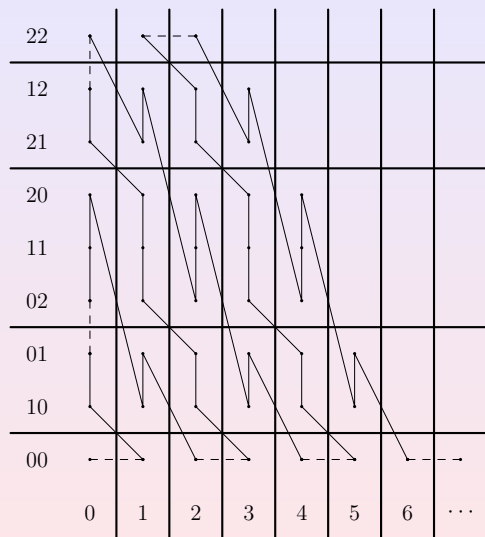
# Bidirectional leveled enumerator for $T^P$

$$\text{BL}(E^1, \text{BL}(E^2, \text{BL}(\dots, \text{BL}(E^P, \text{Null}))))$$



Proof by **induction** that it is a leveled enumerator and a 2-ordering.

## Example of a bidirectional leveled product



The leveled 2-ordering of  $\mathbb{N} \times [0, 2] \times [0, 2]$

# Conclusion and Perspectives

- ▶ Recursive use of bidirectional leveled product gives a **fair** and **2-ordering** of the cartesian product  $T^P$ .
- ▶ Need for a non-recursive version to avoid stack exhaustion
- ▶ Enum package (2300 lines)
- ▶ Self-contained but not yet available separately

Thank you for your attention !